

An Experimental Study of Grouping Crossover Operators for the Bin Packing Problem

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Abstract. The one-dimensional Bin Packing Problem (1D-BPP) is a classical NP-hard problem in combinatorial optimization with an extensive number of industrial and logistic applications, considered intractable because it demands a significant amount of resources for its solution. The Grouping Genetic Algorithm with Controlled Gene Transmission (GGA-CGT) is one of the best state-of-the-art algorithms for 1D-BPP. This article aims to highlight the impact that the crossover operator itself can have on the final performance of the GGA-CGT. We present a comparative experimental study of four state-of-the-art crossover operators for 1D-BPP: Uniform, Exon Shuffling, Greedy Partition and Gene-level; this is the first time that the Uniform, Exon Shuffling and Greedy Partition operators are adapted and studied as a part of the GGA-CGT; moreover, the Uniform crossover has never been used before for solving the 1D-BPP. We measure the performance of the GGA-CGT by replacing its original crossover operator (Gene-level) with each of the other three state-of-the-art operators. Furthermore, we propose a new version of the Uniform crossover and examine two replacement strategies for the Gene-level crossover. Experimental results indicate that the Gene-level crossover operator is shown to have a greater impact in terms of the number of optimal solutions found, outperforming the other operators for the class of Hard28 instances, which has shown the greatest degree of difficulty for 1D-BPP algorithms.

Keywords. Bin packing problem, group oriented crossover operators, evolutionary computation, grouping genetic algorithm.

1 Introduction

The off-line one-dimensional Bin Packing Problem (1D-BPP) is a well-known grouping optimization problem with many applications in logistics, industry, telecommunications, transports, among several others. Given an unlimited number of bins with a fixed capacity $c > 0$ each, and a set of n items, each one with a specific weight $0 < w_i \leq c$, 1D-BPP consists of storing all of the items into the minimum number of bins without exceeding the capacity of any bin. 1D-BPP belongs to the NP-hard class, i.e., the problem complexity grows exponentially as the problem size increases.

It implies that there is no efficient algorithm to find an optimal solution for every instance of 1D-BPP. Searching for the best possible solutions to 1D-BPP, a wide variety of algorithms have been designed. The proposals range from simple heuristics to hybrid strategies, including branch and bound techniques [7], metaheuristics [20] and special neighbourhood searches [4].

However, despite the efforts of the scientific community to develop new strategies, there is not yet an efficient algorithm capable of finding the best solution for all possible 1D-BPP instances, so it is then important to try to identify the characteristics that define the behavior of the algorithms to understand and improve their performance.

One of the suggested methods to solve BBP is the Grouping Genetic Algorithm (GGA) proposed by Falkenauer in 1996 [10], who presented

a design of three new components: (1) a representation scheme for solutions in which groups are seen as genes; (2) a fitness function that evaluates the exploitation of bins' capacity; and, (3) grouping variation operators to modify and re-combine the group-based solutions, including a Segment-level crossover. Later, in 2015, Quiroz-Castellanos et al. [20] proposed the algorithm known as Grouping Genetic Algorithm with Controlled Gene Transmission (GGA-CGT).

Unlike Falkenauer, Quiroz-Castellanos et al. proposed the application of the variation operators in a controlled way, inducing the fullest-bin pattern. GGA-CGT is one of the best algorithms found in the state-of-the-art to solve 1D-BPP; it focuses on the transmission of the best genes on the chromosomes (the fullest-bin pattern), keeping a balance between selective pressure and diversity in the population, in order to favor the generation and evolution of high-quality solutions. GGA-CGT includes an intelligent grouping crossover operator, called Gene-level crossover, which gives the best genes (the fullest bins) a higher probability of being preserved.

The experimental results presented by Quiroz-Castellanos et al. [20] exposed that the Gene-level crossover showed an effectiveness improvement of 30% when compared with the Segment-level crossover proposed by Falkenauer [10]. Despite the success of the Gene-level crossover, the performance of the GGA-CGT is related mainly to the mutation operator, which alone is capable of finding quality solutions.

In the present work, three well-known grouping crossover operators are implemented to solve 1D-BPP: Uniform crossover, Exon Shuffling crossover and Greedy Partition crossover. These operators have never been implemented within the GGA-CGT algorithm before. Furthermore, the Uniform crossover has not been used to solve the 1D-BPP. The goal of the implementation is to measure the performance of these crossover operators with respect to the predefined Gene-level crossover operator. The performance of the GGA-CGT is studied by replacing the original Gene-level crossover operator with each of these state-of-the-art crossovers as well as new

versions of the Gene-level and the Uniform crossover operators.

The paper structure is as follows. Section 2 presents the most relevant state-of-the-art algorithms for 1D-BPP; Section 3 comprises a brief definition of the components of the GGA-CGT; Section 4 includes an explanation for the state-of-the-art grouping crossover operators that will be implemented afterwards; Section 5 contains the experimental proposal to analyze the performance of the mentioned grouping crossover operators; finally, Section 6 summarizes the conclusions and future research paths.

2 Related Work

In the last three decades, different techniques have been implemented in order to find the best solution for the BPP, as it is one of the most interesting problems of the optimization field. Among the techniques that stand out the most are hybrid algorithms and heuristics. For the 1D-BPP study, most algorithms proposed in the literature have been evaluated using a well-studied trial benchmark [8]; it includes 1615 instances in which the number of items n varies within $[50, 1000]$, the bin capacity c is within $[100, 100000]$ and the ranges of the weights are within $(0, c]$.

The specialized literature includes approximation algorithms, which had their performances mathematically analyzed, being the most successful ones: (1) First-Fit Decreasing (FFD), (2) Best-Fit Decreasing (BFD) and (3) Minimum Bin Slack (MBS) [13]. The proposals also include many exact algorithms using dynamic programming, LP relaxation, branch-and-bound, branch-and-price and constraint programming methods [7, 17]. The most relevant results have been obtained by means of metaheuristic and hybrid algorithms covering proposals based on local search [2, 4], evolutionary algorithms [3, 10, 20, 15, 5] and swarm intelligence algorithms [1, 16, 18, 11]. The most exploited approaches, which have allowed obtaining the best results, consist mainly of: (1) the use of simple 1D-BPP heuristics; (2) the application of search space reduction methods; (3) the inclusion of local search techniques based on the dominance criterion; (4) the use of lower

bounding strategies; and (5) the induction of the fullest-bin pattern.

The review of the results obtained by the best 1D-BPP solution algorithms revealed that there are still instances of the literature that present a high degree of difficulty and the strategies included in the procedures do not seem to lead to better solutions. After the literature analysis, it was observed that none of the state-of-the-art strategies had been analyzed to explain the reason for its high-grade or poor performance. Few studies have centered on the analysis of the relationships between the effectiveness of the algorithms and the structure and complexity of the 1D-BPP instances [6]. It is important to understand the algorithms' behavior and evaluate the strategies that allow them to achieve their performance. This work aims at studying the performance of different grouping crossover operators trying to identify the strategies that they use and that positively impact their performance.

3 Grouping Genetic Algorithm with Controlled Gene Transmission (GGA-CGT)

The GGA-CGT proposed in 2015 by Quiroz-Castellanos et al. [20] uses a group-based representation in which each gene represents a group of items or bin. The GGA-CGT is aimed at maximizing the fitness of the individuals in the population. The fitness function is described as follows:

$$F_{BPP} = \frac{\sum_{i=1}^m (S_i/c)^2}{m}, \quad (1)$$

where m is the number of bins in the solution, S_i is the total weight of the items in the bin i and c corresponds to the capacity of the bins. The GGA-CGT algorithm generates an initial population using the FF- n heuristic, in which the n objects of weight greater than 50 percent of the bin capacity are packed in n separate bins, then the remaining objects are accommodated using the well-known First Fit heuristic on a random permutation of this subset.

GGA-CGT uses a controlled selection regarding the choice of individuals to cross and mutate. The strategy consist of an elitist approach together with two inverted rankings to give all the solutions a chance to contribute to the next generation but forcing the survival of the best solutions. For the crossover, n_c parents are selected to generate n_c children. Two sets of parents G and R are generated each with $n_c/2$ individuals, one set being randomly selected from the best n_c individuals with uniform probability (B) and the other set randomly selected from the whole population without considering the elite solutions with uniform probability (R). For mutation, n_m individuals are taken from the best individuals in the population.

GGA-CGT includes a new crossover operator that is referred to as Gene-level crossover, which generates two children c_1 y c_2 from two parents p_1 y p_2 . In this operator, both parents are first sorted in descending order with respect to how full each gene (bin) is. Then, the genes of both parents are compared in parallel, whereby the fuller gene is inherited first. If both genes are equally full, then, for the first child, preference is given to the first parent's gene and, for the second child, preference is given to the second parent's gene. If any parent has more genes than the other, these are inherited directly from this solution. Genes with repeated items are eliminated from the children, and the missed items are reinserted with the FFD heuristic.

Regarding the mutation, it consists of an Elimination operator which works at the gene level, promoting the transmission of the best genes on the chromosome. The Adaptive mutation operator, which considers the bins in descending order of their filling, eliminating the n_b least full bins of the solution and reinserting their items with the Rearrangement by Pairs heuristic. The number of bins n_b to be eliminated from the individual, unlike traditional methods, is calculated in relation to the size of the solution and the number of incomplete bins. The equation is defined below:

$$n_b = \lceil \iota \cdot \epsilon \cdot p_\epsilon \rceil, \quad (2)$$

where ι corresponds to the number of incomplete bins in the solution, ϵ corresponds to the elimination proportion defined by Eq. 3, p_ϵ is the

elimination probability defined by Eq. 4, and k is a parameter that defines the rate of change of ϵ and p_ϵ with respect to i ($k > 0$):

$$\epsilon = \frac{(2 - (i/m))}{i^{(1/k)}}, \quad (3)$$

$$p_\epsilon = 1 - \text{uniform}(0, \frac{1}{i^{1/k}})). \quad (4)$$

The replacement strategy preserves the population diversity and the best solutions by replacing duplicated fitness individuals and the worst fitness solutions with new offspring. The controlled replacement method for the crossover consists of introducing the n_c children such that $n_c/2$ replace the individuals in the set of random parents R and the other $n_c/2$ replace the individuals with repeated fitness first, if there are still un-reinserted children and no solutions with repeated fitness, they are added by replacing the solutions with the worst fitness solutions.

When the mutation operator is applied, some of the elite solutions whose age is less than a predefined *life_span* parameter are cloned. Every clone can be entered into the population in two ways; first, by replacing solutions with repeated fitness, then if there are no solutions with repeated fitness, they are added by replacing the worst fitness solutions.

The details of the heuristics used to generate the population, the rearrangement heuristics to repair solutions, as well as the remaining mechanisms and the parameter settings can be consulted in the work of Quiroz-Castellanos et al. [20].

In order to identify the impact of the Gene-level crossover operator (GLX) on the GGA-CGT performance, an experimental study was performed by using nine different values for the crossover rate, to vary the number of individuals selected for the crossover process (n_c). Fig. 1 and Fig. 2 present the number of optimal solutions found and the average number of generations with different configurations for the GLX. The figures allow observing how the crossover operator seems to have a low impact on the performance of the GGA-CGT. As it can be seen from figures the inclusion of the GLX operator improves the performance of the GGA-CGT, however an

increase in the crossover rate does not seem to contribute to the effectiveness of the algorithm.

The following sections will present a series of studies consisting of: (1) implementing the state-of-the-art crossover operators in the GGA-CGT algorithm; (2) performing experimentation with different crossover percentages within the algorithm; and (3) analyzing the results of the GGA-CGT algorithm with each operator.

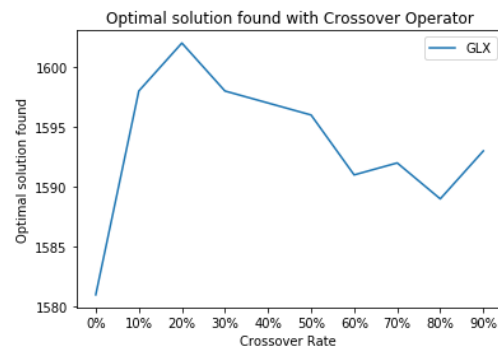


Fig. 1. Number of optimal solutions obtained by the GGA-CGT with the original Gene-level crossover (GLX) for different crossover rates

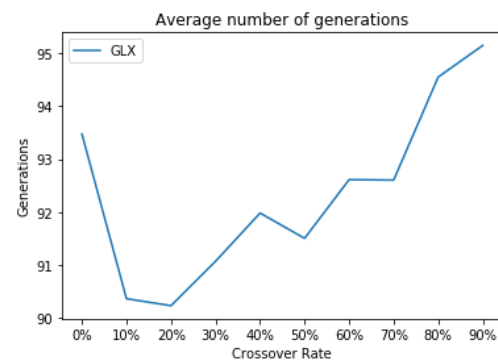


Fig. 2. Average number of generations executed by the GGA-CGT with the original Gene-level crossover (GLX) for different crossover rates

4 Grouping Crossover Operators

The crossover is one the most frequently used genetic operators. This operator combines the

information of two or more solutions (called parents) to produce descendants (called children or offspring). The state-of-the-art algorithms includes different grouping crossover operators (see Ramos-Figueroa et al. [21] for a survey of grouping genetic operators). Grouping crossover operators perform the transmission of the genetic material considering the characteristics of every gene, performing a more controlled combination process, by giving the best genes a higher probability of being preserved.

The state-of-the-art algorithms present four grouping crossover operators that work with the group-based encoding: Gene-level crossover (GLX), Uniform crossover (UX), Exon Shuffling crossover (ESX), and Greedy Partition crossover (GPX). GLX was described in Section 3. The following sections describe the general procedure of the other three operators.

4.1 Uniform Crossover

The Uniform crossover (UX) variation operator generates two children c_1 and c_2 by combining the genes of the parents solutions p_1 and p_2 . In this operator the parents are considered in parallel, i.e. gene 1 of p_1 taken on a par with gene 1 of p_2 . For every pair of genes of the parents, a random value with uniform distribution in the interval (0,1) is generated; a value less than or equal to 0.5 indicates that child c_1 and c_2 receive the gene from the parent p_1 and p_2 respectively, otherwise child c_1 receives the gene from p_2 and offspring c_2 receives the gene from p_1 . During this process, there is the possibility of creating infeasible solutions, which is why heuristics are used to repair them. The performance of the UX operator has only been tested in a Grouping Genetic Algorithm for the Multiple Knapsack problem in 2008 [12]. In our implementation of UX for 1D-BPP, the genes with repeated items are eliminated from the children, and the missed items are reinserted with the FFD packing heuristic. Figure 3a depicts an example of the crossover process followed by UX.

4.2 Exon Shuffling Crossover

The Exon Shuffling crossover (ESX), was proposed by Kolkman and Stemmer (2001) [14]. It has been often used to tackle 1D-BPP [9, 25, 19]. This is an operator that generates a single child c from two parent solutions p_1 and p_2 . The first step consists in joining the parents. Then their genes are ordered from best to worst with respect to their fullness. Finally, the genes are inherited to the child as long as none of the items of the respective gene exist previously in the child [21]. In our implementation of ESX for 1D-BPP, the missed items are reinserted with the FFD packing heuristic. Figure 3b depicts an example of the crossover process followed by ESX.

4.3 Greedy Partition Crossover

The Greedy Partition crossover (GPX), is also among the state-of-the-art grouping oriented crossover operators that have been used to tackle 1D-BPP in other Grouping Genetic Algorithm [23]. This particular operator has two versions and this paper focuses on the group oriented one. Given two parent solutions p_1 and p_2 , GPX generates two children c_1 and c_2 using a greedy heuristic. Here the first step is to order the genes of the parent solutions p_1 and p_2 from most to the least filled. Then, for each child, a vector of probabilities with uniform distribution of the size of the parent solution with more genes is generated. The probability defines from which parent the offspring will receive the gene; if the value generated is less than or equal to 0.5, it indicates that the child c_1 receives the gene from p_1 ; otherwise, it will receive the one from p_2 . The same process is employed to create the child c_2 . Like in the previous cases, in our implementation of GPX for 1D-BPP, the genes with repeated items are eliminated from the children and the missed items are reinserted with the FFD packing heuristic. Figure 4c depicts an example of the crossover process followed by GPX.

4.4 Gene-Level Crossover

The operator is described in Section 3. This operator was proposed by Quiroz-Castellanos et al. [20] for the GGA-CGT algorithm and has been

used several times to solve the 1D-BPP due to its performance [15, 22, 24]. Figure 4d depicts an example of the crossover process followed by GLX.

Since the above operators, are considered the best grouping crossover operators amongst the state-of-the-art, the next section comprises an analysis of these operators to determine which one enables the GGA-CGT to reach the best performance for 1D-BPP. The experiments cover: (1) the integration of the operators in the GGA-CGT algorithm; (2) the study of different crossover percentages; and, (3) the robustness analysis of the results of the best crossover operators inside the GGA-CGT algorithm.

5 Experimentation and Results

This section presents the experiments to analyze the way the different crossover operators can impact on the performance of GGA-CGT. The experimental design consists of three phases. The first one covers the analysis of the state-of-the-art grouping crossover operators (UX, ESX, GPX and GLX) to determine which ones have the best impact on the effectiveness of GGA-CGT. The second one comprises an analysis of the best operators to observe the influence of the number of children generated and their reinsertion to the population. Finally, the third one studies the robustness of the GGA-CGT with the best crossovers, comparing with the original GLX operator.

The performance assessment of each operator involves solving the 1615 standard instances [8], which are distributed among nine sets: data set 1 (720 instances), data set 2 (480 instances), data set 3 (10 instances), triplets (80 instances), uniform (80 instances), hard28 (28 instances), was 1 (100 instances), was 2 (100 instances), and gau 1 (17 instances). Sets data set 1, data set 3, gau 1 and hard28, have shown to have test cases with a high degree of difficulty. Standing out hard28, where there is a higher number of instances that the algorithms cannot solve optimally.

To analyze the performance of each operator on the GGA-CGT, for each instance, a single execution of the algorithm GGA-CGT was run, with the initial seed for the random number generation

set to 1. For each operator ten different crossover rates were explored, from 0% to 90% of the population. For all the parameters, different from the crossover rate, we used the configuration proposed by Quiroz-Castellanos et al. [20].

5.1 State-of-the-Art Grouping Crossover Operators

The GGA-CGT employs the controlled reproduction technique proposed by Quiroz-Castellanos et al. [20], for the state-of-the-art crossover operators we used the same strategy. Both, the UX operator and the GPX operator, generate a set C of n_c children from n_c parents. Like in the original GGA-CGT, the first $n_c/2$ children are introduced to the population replacing the individuals in the set of random parents R . The other $n_c/2$ children are introduced replacing individuals with repeated fitness and replacing the worst solutions. On the other hand, concerning the ESX operator, where $n_c/2$ children are generated from n_c parents (since only one child is generated for every two parents) the reintegration into the population is done in one way: the $n_c/2$ children are introduced replacing the individuals in the set of random parents R .

Table 1, Table 2, Table 3, and Table 4 show the results obtained by the GGA-CGT with each of the state-of-the-art crossover operator (UX, ESX, GPX and GLX, respectively). For every class of instances, each table first shows the number of test cases (Inst.), followed by the number of optimal solutions found by the GGA-CGT with each crossover rate. The last row of each table shows the total number of optimal solutions obtained by each configuration. Moreover, Fig. 5 and Fig. 6 show the number of optimal solutions that are found and the average number of generations executed, when we explore ten different crossover rates (from 0% to 90%) over the four different crossover operators.

The results obtained from implementing the UX variation operator in the GGA-CGT algorithm, are shown in Table 1. As can be seen, the highest number of instances solved is 1592, with a crossover rate of 30%, however, with 90% of crossover rate, the performance of the GGA-CGT is affected since it only solves 1574 of the

Consider a 1D-BPP instance with bin capacity equal to 1000 and 18 items $N = \{1, \dots, 18\}$ with weights:

Item	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
weight	366	268	366	430	263	307	414	287	299	495	251	254	474	252	274	370	269	361

For each gene in the parents a random value is generated using a uniform distribution:

parent 1 : items	1,3	9,15,18	2,4,5	7,10	6,8,11	12,14,17	13,16
fullness	732	934	961	909	845	775	844

parent 2: items	1,17,18	10,13	4,7	2,3,5	6,9,11	8,12,14	15,16
fullness	996	969	844	897	857	793	644

a) UX

For each gene in the parents a random value is generated using a uniform distribution:

0.66	0.21	0.49	0.42	0.06	0.73	0.26
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Offspring 1 items	1,17,18	9,15,18	2,4,5	7,10	6,8,11	8,12,14	13,16
fullness	996	934	961	909	845	793	844

Offspring 2 items	1,3	10,13	4,7	2,3,5	6,9,11	12,14,17	15,16
fullness	732	969	844	897	857	775	644

Free items	Offspring 1	3,9,12,14,15	Offspring 2	2,5,8
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To complete the offsprings, we use the FFD packing heuristic to reinsert again the free items and obtain the new solutions:

offspring 1-r items	1,17,18	2,4,5	7,10	6,8,11	13,16	3,9,12	14,15
fullness	996	961	909	845	844	919	526

offspring 2-r items	1,2,3	10,13	4,7	6,9,11	12,14,17	15,16,5	8
fullness	1000	969	844	857	775	907	287

b) ESX

The genes of both parents are merged and considered in descending order of the fullness:

Offspring items	1,17,18	10,13	2,4,5	6,9,11	8,12,14	15,16
fullness	996	969	961	857	793	644

Free items	Offspring	3,7
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To complete the offspring, we use the FFD packing heuristic to reinsert again the free items and obtain a new solution:

Offspring items	1,17,18	10,13	2,4,5	6,9,11	8,12,14	15,16	3,7
fullness	996	969	961	857	793	644	780

Fig. 3. An example of each state-of-the-art crossover operator for a 1D-BPP instance: UX (Uniform Crossover), ESX (Exon shuffling Crossover), GPX (Greedy partition Crossover), and GLX (Gene-level Crossover) (Part 1)

c) GPX

The genes of each father are considered in descending order of their fullness and a probability vector is generated for each child:

p1	0.68	0.24	0.03	0.99	0.57	0.44	0.76
p2	0.96	0.69	0.50	0.69	0.05	0.59	0.45

offspring 1	items	1,17,18	9,15,18	7,10	6,9,11	4,7	12,14,17	15,16
	fullness	996	884	909	857	884	775	644

offspring 2	items	1,17,18	10,13	2,3,5	6,9,11	13,16	8,12,14	1,5
	fullness	996	969	897	857	884	793	772

Free items offspring 1 2,3,4,5,8,12,13,14 offspring 2 4,7,15,16

To complete the offsprings, we use the FFD packing heuristic to reinsert again the free items and obtain the new solutions:

Offspring 1-r	items	1,17,18	7,10	6,9,11	2,15,16	3,4	5,8,12	13,14
	fullness	996	909	857	912	796	804	726

Offspring 2-r	items	1,17,18	10,13	2,3,5	6,9,11	8,12,14	4,7	15,16
	fullness	996	969	897	857	793	844	644

d) GLX

The genes of each father are considered in descending order of their fullness:

offspring 1: items	1,17,18	2,4,5	9,15,18	7,10	2,3,5	6,9,11	6,8,11	13,16	4,7	8,12,14	12,14,17	1,5	15,16
fullness	996	961	884	909	897	857	845	844	844	793	775	782	644
offspring 2: items	1,17,18	2,4,5	9,15,18	7,10	2,3,5	6,9,11	6,8,11	4,7	13,16	8,12,14	12,14,17	1,5	15,16
fullness	996	961	904	909	897	857	845	844	844	793	775	772	644

Free items Offspring 1 3,15 Offspring 2 3,15

To complete the offsprings, we use the FFD packing heuristic to reinsert again the free items and obtain the new solutions:

Offspring 1-r: items	1,17,18	2,4,5	7,10	6,9,11	13,16	8,12,14	3,15
fullness	996	961	909	857	844	793	640

Offspring 2-r: items	1,17,18	2,4,5	7,10	6,9,11	13,16	8,12,14	3,15
fullness	996	961	909	857	844	793	640

Fig. 4. An example of each state-of-the-art crossover operator for a 1D-BPP instance: UX (Uniform Crossover), ESX (Exon shuffling Crossover), GPX (Greedy partition Crossover), and GLX (Gene-level Crossover) (Part 2)

1615 instances less than those solved with 0% crossover rate. In terms of the average number of generations, with a 90% of crossover rate, the algorithm iterates 100.22 generations per instance, as it is shown in Fig. 6.

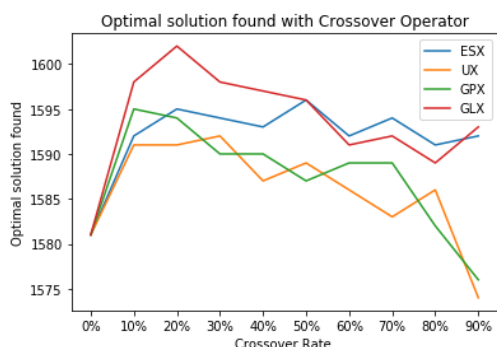


Fig. 5. Number of optimal solutions obtained by the GGA-CGT with each state-of-the-art crossover for different crossover rates

The results obtained from implementing the ESX variation operator in the GGA-CGT algorithm are shown in Table 2. As can be seen, the highest number of the optimally solved instances is 1596, with a 50% crossover rate. Something interesting in this operator is that the optimally solved instances are between 1592 and 1596, and it is not affected with any crossover rate. Concerning to the average number of generations, the algorithm performs a maximum of 91.6 generations, corresponding to a crossover rate of 80%, as it is shown in Fig. 6.

The results obtained from implementing the GPX operator in the GGA-CGT algorithm, are shown in Table 3. The operator solves optimally at most 1595 instances with a 10% crossover rate. As the UX operator, the performance of the GGA-CGT is affected with a 90% crossover rate, since it only finds the optimal solution of 1576 instances less than those solved with 0% crossover rate. It is important to mention that it performs 99.56 generations on average with a 90% crossover rate, as it is shown in Fig. 6.

The same experiment was performed for the crossover operator of the GGA-CGT algorithm, GLX [20]. The results are presented in the Table 4, which with a 20% crossover rate solves optimally

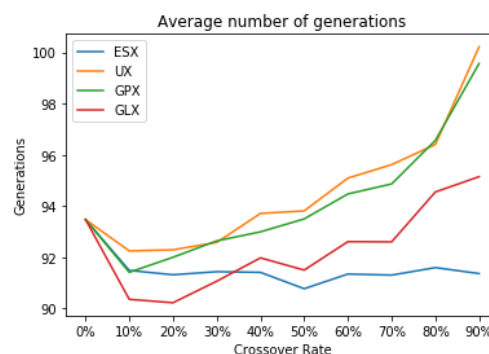


Fig. 6. Average number of generations executed by the GGA-CGT with each state-of-the-art crossover for different crossover rates

1602 instances out of 1615. The minimum number of optimal solutions is found with a crossover rate of 80%, being 1589, and it indeed benefits the algorithm, since without crossover, it finds only 1581 optimal solutions.

From Fig. 5 and Fig. 6 it can be observed that ESX and GLX outperformed the performance of the other two crossover operators, presenting a more stable behavior with different crossover rates.

After the execution of the GGA-CGT with the four crossover operators it was concluded that the ESX operator with a configuration of a 50% of crossover rate and the GLX operator with a 20% crossover rate, found a higher number of optimal solutions, which is why a series of experiments were performed, based on these results and they are described below.

5.2 Reinsertion of Children to the Population

Considering that the GLX algorithm generates two offspring for each pair of crossed parents, for the following experiments, the replacement criteria within the GGA-CGT algorithm were considered. Two versions of the GGA-CGT algorithm were implemented, where the GLX crossover operator generates only one child for each pair of parents.

For the implementation of the following experimentation, the following three cases were addressed: (1) when the ESX algorithm reinserts offspring into the population by replacing individuals with repeated fitness and, if there are still

Table 1. Results obtained by the GGA-CGT with the Uniform crossover (UX) for different crossover rates

Class	Inst.	Crossover rate									
		0%	10%	20%	30%	40%	50%	60%	70%	80%	90%
data set 1	720	708	718	719	719	718	718	718	718	719	718
data set 2	480	480	480	480	480	480	480	480	479	480	479
data set 3	10	9	9	10	10	10	10	10	9	10	10
triplets	80	80	80	80	80	79	80	78	78	78	69
uniform	80	80	80	80	80	79	80	80	80	79	80
hard28	28	9	11	9	10	8	8	7	6	7	5
was1	100	100	100	100	100	100	100	100	100	100	100
was 2	100	100	100	100	100	100	100	100	100	100	100
gau 1	17	15	13	13	13	13	13	13	13	13	13
Total	1615	1581	1591	1591	1592	1587	1589	1586	1583	1586	1574

Table 2. Results obtained by the GGA-CGT with the Exon Shuffling crossover (ESX) for different crossover rates

Class	Inst.	Crossover rate									
		0%	10%	20%	30%	40%	50%	60%	70%	80%	90%
data set 1	720	708	719	719	719	718	718	719	719	718	720
data set 2	480	480	480	480	480	480	480	480	480	480	480
data set 3	10	9	9	10	9	9	10	9	10	9	9
triplets	80	80	80	80	80	80	80	80	80	80	80
uniform	80	80	80	80	80	80	80	80	80	80	80
hard28	28	9	9	10	10	10	12	9	10	9	8
was1	100	100	100	100	100	100	100	100	100	100	100
was 2	100	100	100	100	100	100	100	100	100	100	100
gau 1	17	15	15	16	16	16	16	15	15	15	15
Total	1615	1581	1592	1595	1594	1593	1596	1592	1594	1591	1592

offspring without reinsertion, by replacing the worst solutions; (2) when the GLX algorithm reintegrates the created offspring into the population by replacing individuals within the set of random parents R ; and, (3) as in the case of the ESX operator, when the GLX algorithm reinserts the offspring into the population by replacing individuals with repeated fitness and, if there are still offspring without reintegration, by replacing the worst solutions.

The case where ESX reinserts the children into the population by replacing individuals within the set of random parents R was presented in Table 2, Fig. 5 and Fig. 6. These results are used in Fig. 7 and Fig. 8 for the ESX-1 operator.

The results obtained by the ESX variation operator, corresponding to the case when using reinsertion into the population by replacing the solutions within the repeated fitness group and worst solutions (ESX-2), are shown in Table 5. The GGA-CGT algorithm manages to solve optimally a maximum of 1593 instances with crossover rates of 30% and 40%. The lowest number of instances that it solves optimally is 1589 with a 10% crossover rate.

On the other hand, the results of the implementation of the GLX operator with the reinsertion into the population by replacing the parents of the random group (GLX-1) are presented in Table 7. As it can be seen, the operator implemented within

Table 3. Results obtained by the GGA-CGT with the Greedy Partition crossover (GPX) for different crossover rates

Class	Inst.	Crossover rate									
		0%	10%	20%	30%	40%	50%	60%	70%	80%	90%
data set 1	720	708	720	719	718	719	718	719	719	718	718
data set 2	480	480	480	480	480	480	480	480	480	480	478
data set 3	10	9	9	10	10	10	10	10	10	10	9
triplets	80	80	80	79	80	79	80	79	80	75	73
uniform	80	80	80	80	80	80	80	79	80	80	79
hard28	28	9	11	12	8	9	6	8	6	6	6
was1	100	100	100	100	100	100	100	100	100	100	100
was 2	100	100	100	100	100	100	100	100	100	100	100
gau 1	17	15	15	14	14	13	13	14	14	13	13
Total	1615	1581	1595	1594	1590	1590	1587	1589	1589	1582	1576

Table 4. Results obtained by the GGA-CGT with the Gene-level crossover (GLX) for different crossover rates

Class	Inst.	Crossover rate									
		0%	10%	20%	30%	40%	50%	60%	70%	80%	90%
data set 1	720	708	719	720	719	718	718	718	718	718	719
data set 2	480	480	480	480	480	480	480	480	480	480	480
data set 3	10	9	9	10	9	9	9	9	9	8	9
triplets	80	80	80	80	79	79	78	76	77	74	78
uniform	80	80	80	80	80	79	80	79	79	79	80
hard28	28	9	14	16	15	16	15	13	14	14	12
was1	100	100	100	100	100	100	100	100	100	100	100
was 2	100	100	100	100	100	100	100	100	100	100	100
gau 1	17	15	16	16	16	16	16	16	15	16	15
Total	1615	1581	1598	1602	1598	1597	1596	1591	1592	1589	1593

the algorithm finds a maximum number of 1599 optimal solutions with a 20% crossover rate, while the lowest number of optimal solutions is 1591 with a 80% crossover rate.

The last experiments, corresponding to the GLX with the reinsertion of the offspring in the population replacing the solutions with repeated fitness and worst fitness (GLX-2), are shown in Table 6. This operator allows the GGA-CGT to find a maximum of 1598 optimal solutions with crossover rates of 40%, 60% and 80%, while the minimum number of instances optimally solved is 1595 for 70% and 90% crossover rates, without taking into account the results when there is no crossover.

Based on the results discussed earlier in this section, with respect to the implementation of the two versions of the ESX operator, it is concluded that the ESX operator using a reinsertion of the children in the group of random parents (ESX-1), has a better performance, since it solves on average 1593.22 instances in an optimal way out of 1615. Moreover, the highest number of optimal solutions found by the ESX operator are solved with this reinsertion (1596). On the other hand, the other reinsertion, replacing parents with repeated and worst fitness solves an average of 1591.55 instances in an optimal way.

Regarding the results of the implementation of the GLX operator in the version that only

Table 5. Results obtained by the GGA-CGT with the Exon Shuffling crossover with replacement to individuals with duplicated fitness and the worst solutions (ESX-2) for different crossover rates

Class	Inst.	Crossover rate									
		0%	10%	20%	30%	40%	50%	60%	70%	80%	90%
data set 1	720	708	718	719	719	718	718	718	718	718	718
data set 2	480	480	480	480	480	480	480	480	480	480	480
data set 3	10	9	9	9	10	10	9	10	9	9	10
triplets	80	80	80	80	80	80	79	79	79	80	79
uniform	80	80	80	80	80	80	80	80	80	80	80
hard28	28	9	8	9	9	10	9	10	10	10	9
was1	100	100	100	100	100	100	100	100	100	100	100
was 2	100	100	100	100	100	100	100	100	100	100	100
gau 1	17	15	14	15	15	15	16	15	15	15	15
Total	1615	1581	1589	1592	1593	1593	1591	1592	1591	1592	1591

Table 6. Results obtained by the GGA-CGT with the Gene-level crossover with replacement to individuals with duplicated fitness and the worst solutions (GLX-2) for different crossover rates

Class	Inst.	Crossover rate									
		0%	10%	20%	30%	40%	50%	60%	70%	80%	90%
data set 1	720	708	719	719	719	718	719	719	719	720	719
data set 2	480	480	480	480	480	480	480	480	480	480	480
data set 3	10	9	10	9	9	10	10	10	9	10	9
triplets	80	80	80	80	80	80	80	80	80	80	79
uniform	80	80	80	80	80	80	80	80	80	80	80
hard28	28	9	12	13	13	14	12	13	11	12	13
was1	100	100	100	100	100	100	100	100	100	100	100
was 2	100	100	100	100	100	100	100	100	100	100	100
gau 1	17	15	15	16	16	16	16	16	16	16	15
Total	1615	1581	1596	1597	1597	1598	1597	1598	1595	1598	1595

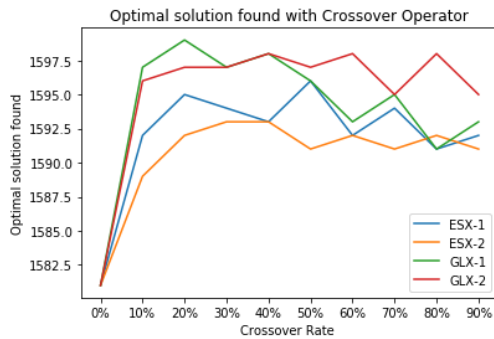
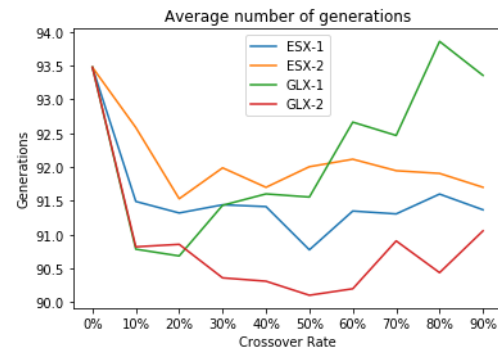
creates one child for every two parents, when the reinsertion into the population is performed by replacing the solutions with repeated and worst fitness (GLX-2), it has better results, finding on average 1596.77 optimal solutions. In the case of the GLX operator with reinsertion to the population by replacing the set of random parents (GLX-1), an average of 1595.44 optimal solutions are found. Moreover, by means of this replacement the highest number of optimal solutions are found by the GLX operator (1599).

5.3 Uniform Crossover with Ordered Genes

After analyzing the different features involved in the crossover operators, we observe that the UX operator was the only one that did not sort the parent solutions in descending order before performing the crossover process. A final implementation of the UX operator was done, where the genes of the parents were first sorted in descending order of their fitness, and then the usual crossover process of the operator was performed (UX-sorted). The results are shown in the Table 8. As we can see, the algorithm with a 10% crossover rate solves a total of 1596 instances in

Table 7. Results obtained by the GGA-CGT with the Gene-level crossover with replacement to individuals in the set of the random parents R (GLX-1) for different crossover rates

Class	Inst.	Crossover rate									
		0%	10%	20%	30%	40%	50%	60%	70%	80%	90%
data set 1	720	708	719	720	719	720	719	718	718	718	719
data set 2	480	480	480	480	480	480	480	480	480	480	480
data set 3	10	9	10	10	9	9	9	10	9	9	9
triplets	80	80	79	79	79	78	79	76	78	75	79
uniform	80	80	80	80	80	80	80	80	79	79	79
hard28	28	9	13	14	14	15	13	14	15	15	11
was1	100	100	100	100	100	100	100	100	100	100	100
was 2	100	100	100	100	100	100	100	100	100	100	100
gau 1	17	15	16	16	16	16	16	15	16	15	16
Total	1615	1581	1597	1599	1597	1598	1596	1593	1595	1591	1593

**Fig. 7.** Number of optimal solutions with different replacements. ESX-1 GLX-1: children replace the individuals in the set of the random parents. ESX-2, GLX-2: children replace the individuals with duplicated fitness and also the worst solutions**Fig. 8.** Average number of generations with different replacements. ESX-1 GLX-1: children replace the individuals in the set of the random parents. ESX-2, GLX-2: children replace the individuals with duplicated fitness and also the worst solutions

an optimal way, while with a 90% crossover rate it affects the performance of the algorithm as it only solves 1574 instances. In addition to the above, a comparison was made between the UX operator and UX-sorted. The results are shown in two tables, one presenting the maximum number of instances solved per class in the (Table 9) and the other with respect to the minimum number of instances solved per class in the (Table 10).

5.4 Robustness of Crossover Operators

To evaluate all of the potential of the best crossover operators, we performed a robustness test by executing three versions of the GGA-CGT thirty times with different seeds of random numbers. These last experiments arose from the comparison of the ESX and GLX operators.

For the GLX operator we used the version in which the children are reinserted in the population by replacing the random parents (GLX-1) with a 50% crossover rate, as well as the original version of the GLX included in the the GGA-CGT

Table 8. Results obtained by the GGA-CGT with the UX-sorted crossover for different crossover rates

Class	Inst.	Crossover rate									
		0%	10%	20%	30%	40%	50%	60%	70%	80%	90%
data set 1	720	708	719	720	719	719	719	719	718	718	718
data set 2	480	480	480	480	480	480	480	479	480	480	478
data set 3	10	9	10	10	10	10	10	10	10	10	9
triplets	80	80	80	80	80	78	79	77	77	76	71
uniform	80	80	80	80	80	80	80	79	79	79	79
hard28	28	9	13	10	8	9	7	7	7	7	6
was1	100	100	100	100	100	100	100	100	100	100	100
was 2	100	100	100	100	100	100	100	100	100	100	100
gau 1	17	15	14	13	13	13	13	13	14	14	14
Total	1615	1581	1596	1593	1590	1589	1588	1584	1585	1584	1575

Table 9. Results of the highest number of optimally solved instances by the GGA-CGT with the UX and UX-sorted crossover operators

Class	Inst.	UX	UX-sorted
		Crossover rate	
		30%	10%
data set 1	720	719	719
data set 2	480	480	480
data set 3	10	10	10
triplets	80	80	80
uniform	80	80	80
hard28	28	10	13
was1	100	100	100
was 2	100	100	100
gau 1	17	13	14
Total	1615	1592	1596

Table 10. Results of the lowest number of optimally solved instances by the GGA-CGT with the UX and UX-sorted crossover operators

Class	Inst.	UX	UX-sorted
		Crossover rate	
		90%	90%
data set 1	720	718	718
data set 2	480	479	478
data set 3	10	10	9
triplets	80	69	71
uniform	80	80	79
hard28	28	5	6
was1	100	100	100
was 2	100	100	100
gau 1	17	13	14
Total	1615	1574	1575

algorithm by Quiroz-Castellanos et al. [20], that generates two children with a 20% crossover rate. On the other hand, for the ESX operator we also used the version in which the children are reinserted in the population by replacing the random parents (ESX-1) with a crossover rate of 50%. These operators were selected as they are within the group of operators that showed the higher effectiveness. The experiment consists of 30 executions of the GGA-CGT algorithm (48,450 runs) with each crossover operator.

The results of the experiments are shown in Table 11, Table 12 and Table 13. The second column (Inst.), corresponds to the total number of instances belonging to the class in the first column. The column "Optimal solutions in the 30 executions" includes the total number of optimal solutions found in the 30 executions.

The column "Average number of optimal" corresponds to the average number of instances optimally solved in each set. Finally, the column

Table 11. Results for the 30 executions of the GGA-CGT with the ESX when it generates only one child

Class	Inst.	Optimal solutions in the 30 executions	Average number of optimal	%
data set 1	720	21562	718.73	99.82%
data set 2	480	14400	480.00	100.00%
data set 3	10	282	9.40	94.00%
triplets	80	2396	79.87	99.83%
uniform	80	2400	80.00	100.00%
hard28	28	299	9.97	35.60%
was 1	100	3000	100.00	100.00%
was 2	100	3000	100.00	100.00%
gau 1	17	467	15.57	91.57%
	1615	47806	1593.53	98.67%

Table 12. Results for the 30 executions of the GGA-CGT with the GLX when it generates only one child

Class	Inst.	Optimal solutions in the 30 executions	Average number of optimal	%
data set 1	720	21562	718.73	99.82%
data set 2	480	14400	480.00	100.00%
data set 3	10	281	9.37	93.67%
triplets	80	2370	79.00	98.75%
uniform	80	2393	79.77	99.71%
hard28	28	394	13.13	46.90%
was 1	100	3000	100.00	100.00%
was 2	100	3000	100.00	100.00%
gau 1	17	476	15.87	93.33%
	1615	47876	1595.87	98.82%

“%” corresponds to the percentage of optimal solutions found.

It is observed that the GGA-CGT algorithm with the GLX operator, which generates two children, found on average a higher number of optimal solutions, with a 98.99% percentage.

While the GLX and the ESX that generate one child found 98.8% and 98.67% of the optimal solutions, respectively. Furthermore, on average, the original GGA-CGT with the GLX operator that generates two children resolves in an optimal way

an average of 1598.67 instances out of 1615. Regarding to the GLX and the ESX that generate one child, they found on average 1595.87 and 1593.53 optimal solutions, respectively.

5.5 Statistical Analysis of the Effectiveness of GGA-CGT with the Best Crossover Operators

A series of statistical tests were applied to the results of the 30 executions of the GGA-CGT algorithm with the ESX, GLX and GLX-V1

Table 13. Results for the 30 executions of the original GGA-CGT with the GLX that generates two children

Class	Inst.	Optimal solutions in the 30 executions	Average number of optimal	%
data set 1	720	21571	719.03	99.87%
data set 2	480	14400	480.00	100.00%
data set 3	10	291	9.70	97.00%
triplets	80	2391	79.70	99.63%
uniform	80	2400	80.00	100.00%
hard28	28	430	14.33	51.19%
was 1	100	3000	100.00	100.00%
was 2	100	3000	100.00	100.00%
gau 1	17	477	15.90	93.53%
	1615	47960	1598.67	98.99%

operators. The statistical test applied was the Wilcoxon Rank-sum test with a 95% confidence. These tests were performed on two result sets. First, the error of the results was calculated for each instance in each execution, which was determined as the relative difference defined as: $(y - x)/x$, where x corresponds to the optimal number of bins and y corresponds to the number of bins obtained.

For the first set of experiments, the average error of the 30 executions was obtained for each instance.

Next, the Wilcoxon test was applied to the average errors of the algorithm per class of instances with the three operators: ESX vs GLX-V1, ESX vs GLX and GLX-V1 vs GLX, obtaining 9 p -values (one for each class). The results are shown in the Table 15.

For the second set of tests, all instances were considered separately. A Wilcoxon test was performed for each instance with the results obtained from the algorithm with the different operators, i.e.: ESX vs GLX-V1, ESX vs GLX and GLX-V1 vs GLX. The Table 16 shows the results of the instances where a significant difference was observed in the results obtained with the different operators.

The Wilcoxon test, at the class level, shows no significant difference between the operators, which

would indicate that the performance is the same among the three. However, the test at the instance level does show a difference, in some instances of the Hard28 and Gau 1 classes, between the ESX operator and the GLX-V1 and GLX operators, indicating that for these instances the ESX operator is the worst operator.

5.6 Difficult 1D-BPP Instances

Finally, with the objective to identify the features of the 1D-BPP instances that present the highest degree of difficulty for the four state-of-the-art grouping crossovers, Table 16 shows the set of instances for which the optimum was not found. These results were obtained from the set of experiments performed with the different crossover rates, selecting those instances that were not solved optimally in at least one of the configurations.

The first column includes the name of the sets of instances that include cases for which the optimal solutions could not been found for some of the configurations of the four crossover operators. The second column contains the name of the difficult instances, and the following columns include an X for each operator that fails to find the optimal solutions for each instance in one of its configurations. The instances belong to the classes: data set 1, data set 2, data set 3, triplets,

Table 14. Instances for which the GGA-CGT does not find the optimal solution with each operator for any of the crossover rates

Class	Inst.	UX	ESX	GPX	GLX	Class	Inst.	UX	ESX	GPX	GLX
data set 1	N3c1w2.r	X				triplets	t501_17				X
	N3c3w4.c	X	X	X	X		t501_18	X			X
	N4c3w4.s	X	X	X	X		t501_19				X
data set 2	N2w1b2r5	X		X		Uniform	u250_12	X		X	X
	N2w1b2r8			X			h1D-BPP832	X		X	X
data set 3	Hard2			X			h1D-BPP40	X	X	X	X
	t60_00	X					h1D-BPP360	X	X	X	
	t60_01	X		X			h1D-BPP645	X	X	X	X
	t60_05	X					h1D-BPP742	X	X	X	
	t60_06	X		X			h1D-BPP766	X	X	X	X
	t60_07	X					h1D-BPP60	X	X	X	X
	t60_08	X					h1D-BPP13	X	X	X	X
	t60_10			X			h1D-BPP195	X	X	X	X
	t60_11	X					h1D-BPP709	X	X	X	X
	t60_12	X					h1D-BPP785	X	X	X	X
	t60_13	X				hard28	h1D-BPP47	X	X	X	X
	t60_14				X		h1D-BPP181	X	X	X	X
	t60_15				X		h1D-BPP485	X	X	X	X
triplets	t60_17	X					h1D-BPP640	X	X	X	X
	t60_18	X					h1D-BPP144	X	X	X	X
	t60_19						h1D-BPP561	X		X	X
	t120_03	X					h1D-BPP781	X	X	X	X
	t120_14	X					h1D-BPP900	X	X	X	X
	t249_04				X		h1D-BPP178	X	X	X	X
	t501_00			X			h1D-BPP419	X	X	X	X
	t501_04				X		h1D-BPP531	X	X	X	
	t501_05	X		X	X		h1D-BPP814	X	X	X	
	t501_08				X		TEST0058	X		X	
	t501_09				X	gau1	TEST0014	X	X	X	X
	t501_14			X	X		TEST0030	X	X	X	X
	t501_16			X			TEST0005	X		X	X
Total								46	25	39	36

uniform, hard28 and gau 1. The highest number of instances where the optimum is not found belong to the triplets class. With respect to the class data set 3 and uniform, only one instance is not optimally solved.

As it can be seen, the instances belonging to the hard28 class are the most complicated for the 4 operators. The operator with the lowest number of optimally solved instances is the ESX

operator. Although the GLX operator optimally solves the highest number of instances, these instances tend to be more variable with different crossover percentages.

The UX operator, on the other hand, is the one that has the greatest diversity with respect to the number of instances not optimally resolved, being a total of 46. With respect to the instances belonging to the triplets class, the ESX operator

Table 15. *p*-value of the Wilcoxon test for the error of the GGA-CGT with the ESX, GLX-V1 and GLX crossovers in the different classes of instances

Class	<i>p</i> -value for the average error		
	ESX vs GLX-V1	ESX vs GLX	GLX-V1 vs GLX
data set 1	0.9637	0.9636	1
data set 2	1	1	1
data set 3	0.9698	0.9698	0.9698
triplets	0.2459	0.1557	0.8618
uniform	0.8914	0.8914	0.9986
hard 28	0.5552	0.5121	0.9151
was 1	1	1	1
was 2	1	1	1
gau 1	0.7828	0.7828	0.9862

is the only operator that has no problem in finding the optimum, besides being the only one to solve optimally all the instances of the Uniform class.

Although the GGA-CGT algorithm finds the optimal solution of most of the well-known benchmark instances within the state-of-the-art, it has been shown that it does not perform well with the 2800 new difficult instances, referred to as $BPPv_{u-c}$, proposed by Carmona-Arroyo et al. [6]. The results that were obtained by the GGA-CGT demonstrated that for most of these instances, the operators included in the GGA-CGT do not appear to lead to better solutions.

6 Conclusion and Future Work

In this research we propose an experimental study about the GGA-CGT for the 1D-BPP, in which different grouping crossover operators were used to compare and measure the performance. The conclusions of this research are presented as follows.

From the experiments carried out in Section 5, it is concluded that the original Gene-level crossover operator (GLX), when generating two children, outperforms the other state-of-the-art grouping crossovers, achieving a better performance with respect to the number of optimal solutions found, being 1602 out of 1615. The second experiment,

also explained in Section 5, on the implementation of the ESX and GLX operators in their two versions (generating only one child, with replacement to random parents, and, with replacement to solutions with repeated fitness and worst solutions), shows that the best operator is the GLX operator with the reinsertion to the population through the replacement to random parents, since it solves optimally a total of 1599 instances out of 1615.

The last experiment allow us to validate that the original GLX operator, that generates two children, is the one that achieves the best performance, showing a robust behavior in different runs of the GGA-CGT. Finally, it can be concluded that the GLX operator performs better within the GGA-CGT algorithm.

Although the GLX operator solves a larger number of instances, per class there is no significant difference between the operators. But there is a difference for some instances of the Hard 28 class, where it has been shown that the GLX operator has the best performance of the operators.

For future work, a detailed review of the operators is planned to identify the strategies that allow for them to improve their performance in the different classes instances, with the objective of using this knowledge to design a new crossover operator with a better performance in terms of time, number of generations and number of optimal solutions found for instances of different classes. It is also intended that the GGA-CGT algorithm with the operator to be designed will improve the performance of the algorithm for the $BPPv_{u-c}$ instances with which it does not perform well.

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Table 16. *p*-value of the Wilcoxon test for the error of the GGA-CGT with the ESX, GLX-V1 and GLX crossovers in the most difficult instances

Class	Instance	<i>p</i> -value of the error of the 30 executions		
		ESX vs GLX-V1	ESX vs GLX	GLX-V1 vs GLX
hard 28	h1D-BPP360.txt	0.0147	0.147	1
	h1D-BPP742.txt	0.0039	0.0009	0.6574
	h1D-BPP47.txt	3.3853x10 ⁻⁷	2.9537x10 ⁻⁸	0.6574
	h1D-BPP640.txt	0.01471	0.01471	1
	h1D-BPP531.txt	2.9537x10 ⁻⁸	2.9537x10 ⁻⁸	1
	h1D-BPP814.txt	0.0039	0.0039	1
gau 1	TEST0030.txt	0.0358	0.069	0.6574

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